

# A Pattern System for Network Management Interfaces

*NMI builders and applications developers should find the operational benefits of a pattern-based framework called Layla a welcomed ally in their efforts.*

**N**etwork management systems are used to control and monitor the components of distributed systems like communication networks, where many different subsystems need to collaborate to offer a service. Communication networks are large dynamic systems, comprising parts from various vendors and evolving over time, with parts getting grafted on and others being removed. Network management is a challenging task that usually requires remote access to widely distributed information coming from various sources. Operations have to be performed on numerous system components. Moreover, the access interface to the components can vary greatly, depending on their nature, type, and manufacturer.

We define a *network management interface* (NMI) as the middle layer of a network management system, situated between the high-level control processes and the low-level components of the system [12]. The lower layers, which usually depend heavily on the execution platform at hand, are not part of the NMI.

Figure 1 illustrates a sample distributed system (network) under management. The system consists of two workstations communicating through a switch. Each workstation has a communication port that is attached to each end of the connection path, and the path itself is contained in the switch. The management system includes a management console that has access and con-

trol over all the components of the network through a symbolic representation provided by the NMI. Note that the NMI must include a number of communication stacks to access all the various components.

International standardization bodies have produced various tools for defining network management systems and their NMIs. Among the most advanced tools are the Common Management Information Service (CMIS) of Open Systems Interconnection (OSI) [12], and the Simple Network Management Protocol (SNMP) of the Internet [9]. Whereas CMIS, along with its protocol for information exchange between systems CMIP, is based on the object-oriented para-

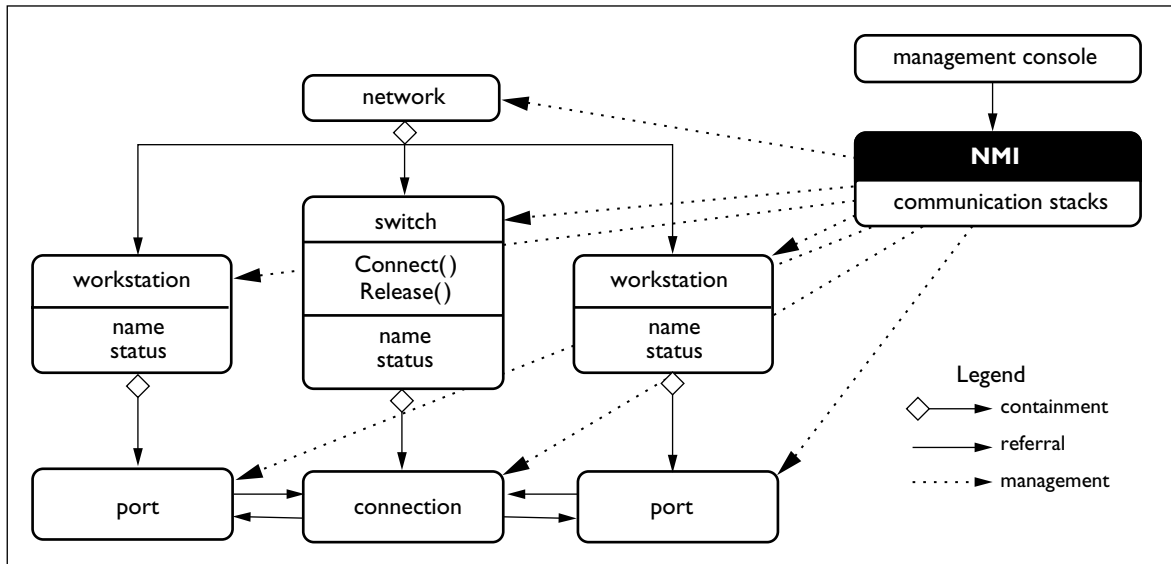


Figure 1. Sample network managed by network management system

digm, SNMP uses tables similar to those used in the relational model of databases. However, SNMP is moving toward the OO paradigm; its new version SNMPv2 embodies some notion of inheritance.

In the IGLOO project, we have developed a prototype pattern-based framework for NMIs called “Layla.”<sup>1</sup> In designing Layla, we wanted to leverage off commercial implementations of standardized network management protocols, and came up with a number of wrapper classes that encapsulate the specific details of any particular protocol engine. Layla supports OSI NMIs and includes provisions for the OO nature of CMIS that are not necessarily found in other protocols for network management. We have built several NMIs to date using Layla, in cooperation with Teleglobe Canada Inc., our main industrial partner.

In Layla’s early development we decided to take an approach based on design patterns. Aware of their potential for making architectures easier to modify,

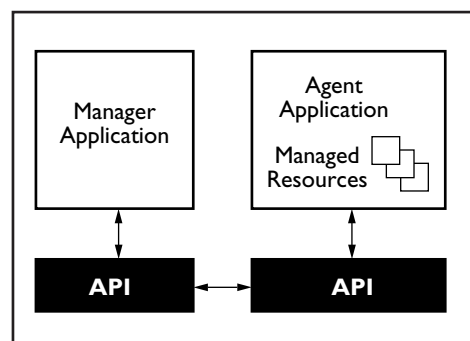
maintain, and reuse, and of their documentation value [2, 3, 7, 11], we wanted to verify whether these qualities also hold in a domain as complex as NMIs, and in which way domain-specific aspects would come into play. A further objective was to organize the system of patterns that would result from our design in a systematic and coherent way. The resulting framework architecture can be described as a heterogeneous system of design patterns. The system consists of previously published, general-purpose patterns, several new and domain-specific patterns taken from NMI standards, as well as a couple of basic patterns relevant in Layla’s application programming interface (API). The patterns are implemented as a system of C++ classes that form the framework. The framework encapsulates to a large extent the underlying communication API, as shown with the adoption of two different, commercially available APIs.

Many software manufacturers offer packaged solutions for implementing NMIs. Such solutions are often called APIs. An API typically includes data structures

<sup>1</sup>In Zohar, Layla is an angel in charge of the newly created spirits.

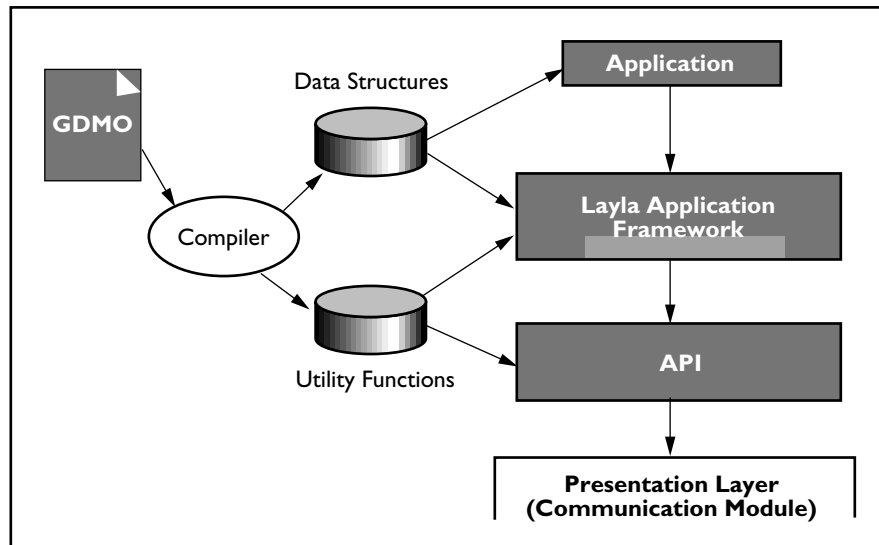
### How it Works

In a typical network management system, the managed resources are grouped under an agent application and controlled remotely by a manager application. The agent and the manager use both an underlying API to communicate with each other and to exchange management information. Commercially available APIs exhibit limitations that may be overcome by encapsulating them in an application framework.

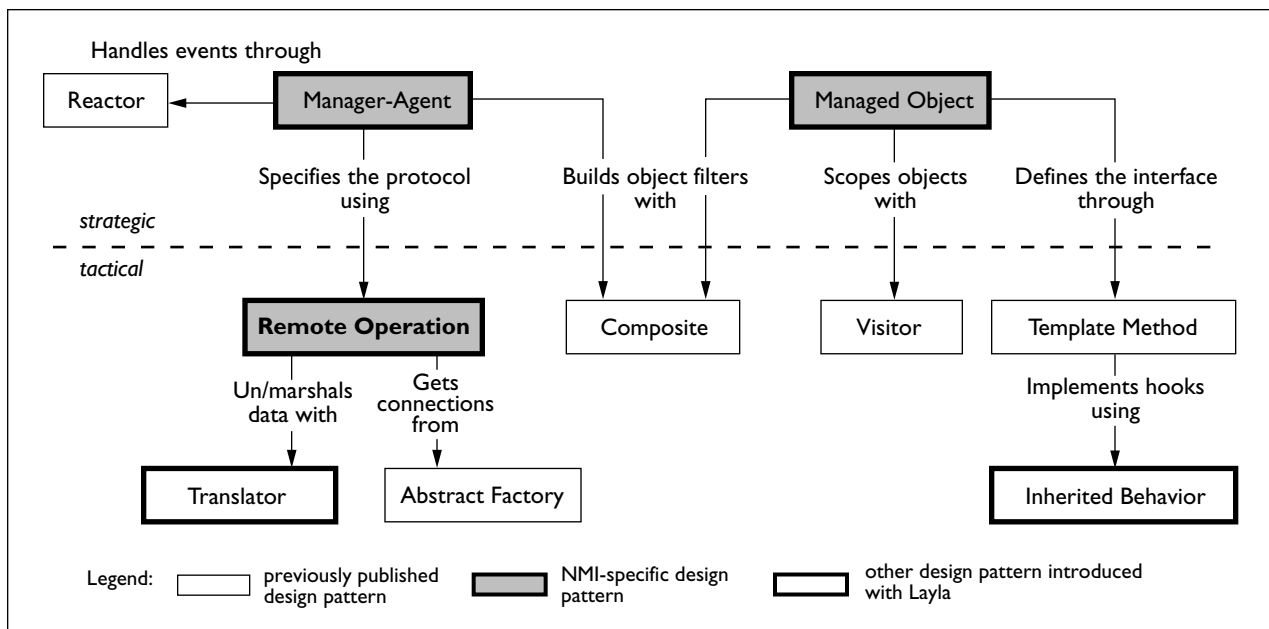


and function prototypes, as well as a set of precompiled libraries that implement those functions. A developer can describe a network management function in terms of API calls and implement the NMI by reusing the code in the libraries (see “How it Works” for more details).

Developing an NMI based on an API usually involves a number of steps (see Figure 2). First, the developer must specify the NMI using the specification language(s) supported by the API. For example, CMIS uses one language for describing object classes and relationships (GDMO) and another for the data structures used by the objects (ASN.1). In contrast, the



**Figure 2.** Layered structure of NMI, with Layla mediating between object-oriented (dark grey) and procedural (light grey) components



**Figure 3.** The pattern system of Layla

SNMP standard uses only one specification language to describe data structures (ASN.1), and the semantics are described using plain English. Once the specification has been written, an automated tool is used to map it to a specific programming language, often C. The tool, typically some sort of compiler, needs to generate the appropriate data structures and utility functions in the target implementation language. Once generated, the developer can use these structures and functions to implement the NMI. The data structures are used by the application to pass information to the API. The utility functions are used by the application to manage those structures, and by the API to transfer

data across process boundaries via the presentation layer (communication module). The API also comprises management protocols for data exchange across the various computing platforms typically found in today’s large and heterogenous networks.

Network management APIs are a powerful development tool, since they take care of many low-level communication issues such as connection establishment/release, buffering, synchronization, and type conversion between different computing platforms. Still, they exhibit several shortcomings. First, most

APIs provide only an interface to the C language, whereas many NMI specification languages are object-oriented, imposing on the developer a discrepancy of paradigms. Second, applications are highly dependent on the output of the specification compiler, and any changes either to the initial specification or to the compilation mechanism may require major modifications to the application code, ultimately leading to serious maintenance problems. Furthermore, many APIs offer functionality that is not necessarily relevant to the NMI under development, yet they are often limited to the lowest common denominator of the functionalities provided by the various flavors of NMIs.

We feel these shortcomings are best addressed by devising an application framework to mediate between the application and the API. The developer then simply has to deal with the OO specification of the NMI and the framework, which itself is object-oriented, yet comprises a procedural kernel for interfacing with the API (Figure 2). Also, a framework encapsulating the underlying API makes NMIs independent of the API that is actually used. Furthermore, generated code can be hidden within the framework, and the framework can be designed such that it provides minimum but sufficient functionality for building NMIs, including high-level services. Note that introducing such an additional framework layer may lead to performance degradation. Our experience in the project, however, suggests that with careful framework design, performance loss can be kept to a minimum.

For our framework development, Teleglobe supplied us with two different APIs, both geared toward OSI NMIs as defined by the CMIS standards [12], and providing the low-level functionality required. One of them, used at the beginning of our project, is Base-WorX from AT&T [1], the other one is DM from Hewlett Packard (on top of which the current version 1.1 of Layla is built [4]). We have strived for complete encapsulation of the API into the framework, so the classes defined for a given NMI can be compiled with one API or the other, thus making the NMI independent of the API that is actually used. However, since the NMI is dependent on the structures generated by the

## Documenting Patterns and Pattern Systems

Patterns capture the essence of successful solutions to recurring problems that arise when building software systems. They are usually documented, especially when grouped together in a catalog, in a common format, or template. A popular format is the one suggested by Gamma et al. [3]. The format consists of: intent, motivation, applicability, structure, participants, collaborations among participants, consequences, implementation suggestions, known uses, and related patterns.

Patterns should be interwoven in pattern systems [2, 8] that describe how they are connected and how they complement each other. For some small and well-known domains, pattern systems have been devised to cover all and every aspect of importance in the domain, and thus could be called “pattern languages.” In most other domains, such as NMIs, patterns only cover certain aspects of software construction. However, as the pattern discipline matures, more pattern languages are likely to appear.

All new patterns found in Layla have been described using the format of Gamma et al. [3], and are available, together with further documentation of Layla, at [www.iro.umontreal.ca/labs/gelo/layla/](http://www.iro.umontreal.ca/labs/gelo/layla/). The individual pattern descriptions comprise numerous hints and references to relevant patterns within and beyond Layla. The classification by their role (tactical or strategic) and scope (NMI-specific or general-purpose), together with the description of the interplay of the patterns in Layla applications, further describes the pattern system.

specification compiler that comes with the API (see Figure 2), this encapsulation can never be complete.

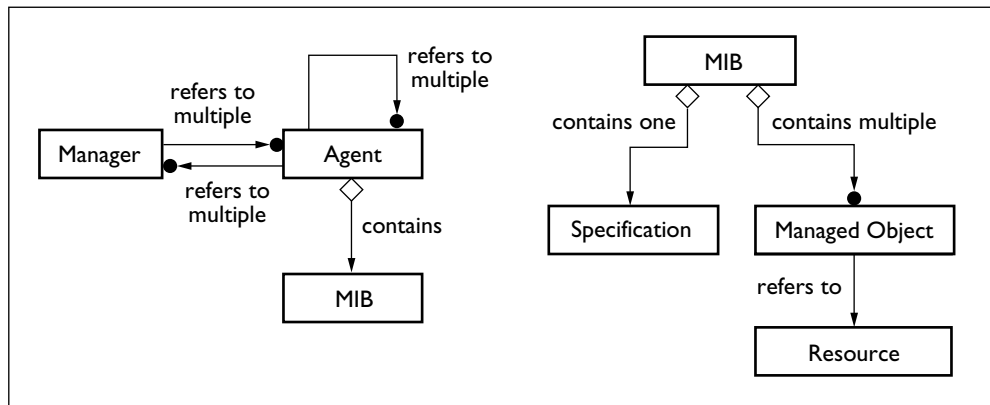
## The Pattern System of Layla

Figure 3 depicts the pattern system that constitutes the architecture of the Layla framework. Strategic patterns [11] are placed at the top, and the more generic but less critical tactical patterns are shown in the region below the separator line. Arrows stand for use-relationships, indicating the originating pattern uses the destination pattern for the functionality mentioned on the label of the arrow. (See the “Documenting Patterns and Pattern Systems” sidebar for details.)

Each major task in Layla is described by a design pattern, many of which come from the literature [3, 10]. This indicates the published design patterns are indeed expressive and generic enough to be easily applied to a new application domain. Several tasks and design solutions that were not design patterns originated from the constraints imposed by the underlying network management standards. Since these solutions are applicable to virtually any network management system, they can be thought of as NMI-specific, hence domain-specific design patterns (Manager-Agent,

Managed Object, Remote Operation). Two further tasks were considered flexible and generic enough to be applicable to other fields, and thus were described as design patterns, too (Translator and Inherited Behavior). They were then circulated on mailing lists and

generated some interesting feedback from the design pattern community. (The interaction of the Layla patterns is illustrated in the accompanying sidebar.)



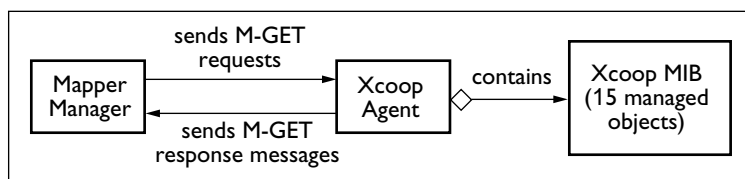
**Figure 4.** Structure of Manager-Agent (left) and Managed Object patterns (right)

### Layla's Patterns at Work

Many network protocols such as Xcoop use network management functions to reserve and activate physical connections. Among the applications developed for validating the Layla framework, we have implemented an Xcoop agent for a transit system. The Xcoop agent and an associated mapper manager interact according to the Manager-Agent design pattern, with events being handled by the Reactor pattern. The agent contains the Xcoop MIB, a collection of 15 managed objects implemented using the Managed Object, Template Method, and Inherited Behavior patterns. The managed objects are grouped using the Compositor pattern. An implementation of the Abstract Factory pattern is used to connect the agent with the manager and to enable CMIP operations (implemented using the Remote Operation and the Translator patterns). The manager sends M-GET requests to the top object in the agent's containment hierarchy with a scope designating the whole tree. When the agent receives such a request, it uses a CScope object (an implementation of the Visitor pattern) to enact the M-GET request in each and every managed object of the MIB. Each managed object gives rise to an M-GET response message back to the mapper manager, which sorts them out and uses that information to display a map of the MIB.

In building this application, the managed objects' specifications must be written in GDMO and ASN.1. For implementing the managed objects, the agent, and the manager, much of the code is supplied by Layla, and only minor portions need to be defined by the developer, mostly through subclassing.

Xcoop is a European ATM Pilot project.



### Manager-Agent Pattern

Imagine a large system of collaborating components (or resources) that provide a service, such as a telecommunication network. Such a system is often managed from a central console that controls all the components in the system. This console is typically called a "manager." It issues commands to monitor and regulate the system, and listens to failure alarms.

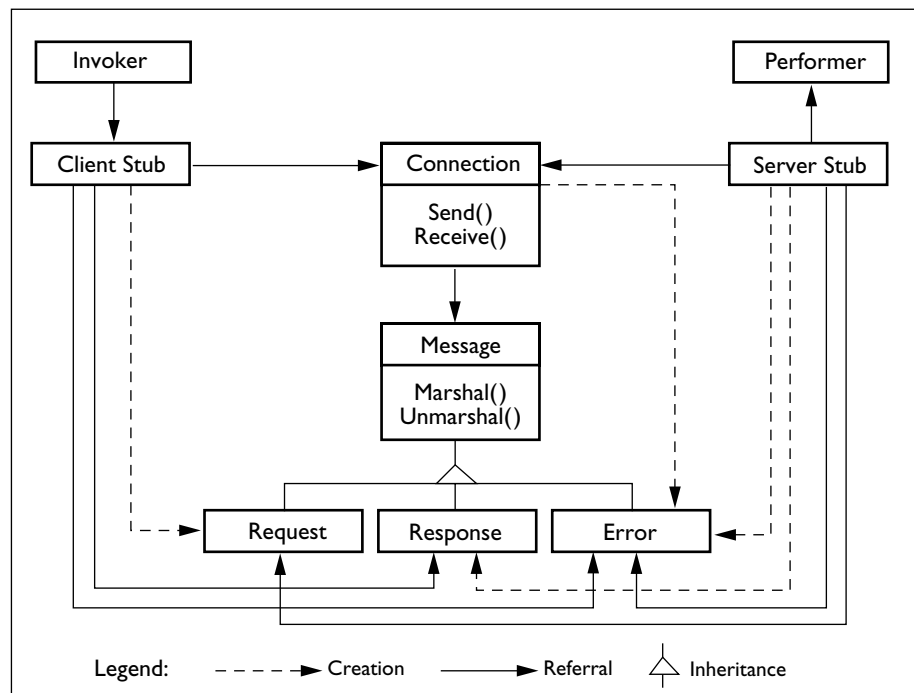
The biggest problem in such a system is to control the complexity of the manager. When attempting to solve this problem, the developer must deal with a number of conflicting forces. For one thing, the manager must be powerful enough to handle its task, and yet present a unified management interface for the whole system as well as for the subsystems thereof. Furthermore, there can be a large variety of management functions to be performed, increasing the complexity of the manager. There can also be an extremely large number of components to be managed, a setting that might overload any single manager. To further complicate matters, the managed components typically come in various kinds of management interfaces, functionalities, and semantics, all of which must be presented through the manager interface. Finally, it might be desirable to have a portion of the system manage itself automatically, so as to relieve the load of the overall manager.

The proposed solution is to first isolate the management functionalities in one or more Manager objects that han-

dle all the management aspects of the system. Then, the set of components is partitioned and brought under the control of individual Agents. The Agents take the responsibility for a group of related resources (functional nature, logical relationship, manufacturer, and so forth). Each Agent represents the whole subsystem under its jurisdiction toward the Managers. Each Manager will interact with multiple Agents within the system in order to handle a given management task. Similarly, each Agent can report to more than one Manager. This solution may be applied recursively, in order to simplify Agents that are in charge of large subsystems. The Agents may then act as the local Managers of the subsystems under their jurisdiction (see Figure 4, left). The set of resources grouped under one Agent is a Management Information Base (MIB).

Applying the Manager-Agent pattern results in a decoupling of the management policy from the system being managed, with the policy being implemented by the set of Managers. Management responsibilities can be spread across multiple Managers, delegated to subsystems (as part of an Agent's functionality), or can be a combination of both. In this way, a divide-and-conquer approach can be applied to managing the system, dividing the whole management task into a set of smaller, more manageable subtasks.

The Manager-Agent pattern brings with it a number of benefits, but also some liabilities. On the positive side, the Manager and the Agent use a single protocol to communicate. This encapsulates the proprietary protocols used by the resources in the Agent and simplifies the implementation of the Manager. The Manager is thus able to communicate with any and all the resources in the system, regardless of their origin or nature. In this capacity, the Agent acts as a large application of the Adapter pattern [3]. Furthermore, the hierarchy of the system is expressed through the organization of the Agents. This means the Manager does not need to maintain its own map of the system, but can rely on the encapsulation provided by the Agents. Modifications in one area of the system need only be reflected in the relevant Agent, minimizing the effect on the Managers that manage that area.



**Figure 5.** Structure of Remote Operation pattern

Finally, the Manager-Agent pattern can be seen as a symmetrical variation of the client/server architecture, where the Agent plays the role of the server and the Manager that of the client. In traditional client/server interactions, all the interactions originate from the client's side, and not the server's side. In the Manager-Agent pattern, however, both agent and manager can be the instigator of an interaction at any time, either via a command going from a Manager to an Agent, or via an alarm going from an Agent to a Manager.

On the downside, each Agent adds a level of indirection when the Manager needs to access the Agent's resources. The Agent, in order to accurately portray the MIB, must either maintain a special internal database or apply a set of translation rules every time a request is made. In both cases, there is an added layer of processing when accessing the managed resources, and this layer can degrade performance when manipulating large resource bases. Whichever mechanism is selected, it needs to be implemented with care. As a positive side effect, though, the Agent can encapsulate vendor-specific details of the resources, furthering interoperability in the system.

Note that the relationships between Managers and Agents must be maintained adequately. One approach is to have each Manager and Agent maintain their own list of collaborating opposites, resulting in a rather inflexible update scheme. Instead, the Mediator [3] pattern may be used to maintain all these relation-

ships. Alternatively, the Remote Operation pattern may be used to provide location transparency. Yet another approach is to use the Broker pattern discussed in [2]. In this article, we will only cover the Remote Operation pattern, the solution we adopted in Layla for its simplicity of implementation; before doing so, we take a closer look at the Managed Object pattern governing the MIBs.

### Managed Object Pattern

You need to manage a large collection of resource entities. Each entity has an interface that may somewhat vary from the interfaces of the other entities, resulting in a great number of different interfaces with somewhat similar features. Such variation can be caused by the nature of the entities, their manufacturers, and so on. In addition, there is often a hierarchical and/or a containment relationship among the entities being managed. The problem here is to provide a unified interface for representing and controlling the resources, while still allowing for a fine degree of control. Although the resources may have diverging interfaces, they need to be controlled in a uniform way, and the relationships among them should be taken into account.

The proposed solution is to use an instance of the Adapter pattern [3] for each individual Resource, in order to translate its particular interface into one shared by all Resource instances. We call this Adapter a Managed Object. It enables management operations to be performed using a single interface. The Managed Object instances can then be aggregated using the Composite pattern [3]. The composition is called a MIB. It is sensitive to all the relationships that usually exist among the Resource instances, such as hierarchy and containment. The layout and interface of the MIB follows a Specification that describes the content of the MIB to the outside world (see Figure 4, right).

The result is a flexible structure that appropriately mirrors the organization of the Resource instances and that uses a single interface for the management of these instances. An individual Resource can be located by traversing the MIB (with the Iterator pattern and/or the Visitor pattern [3]) and by applying operations on the Managed Object's interface. It will in turn apply the corresponding operations to the Resource using the appropriate interface. The Managed Object, representing the Resource being managed, supplants the Resource's specific interface with the interface that is expected by the management system, and it implements management-specific operations such as the `VisitMe()` operation in the Visitor pattern. The MIB acts as a repository for all the management data and functionalities of the system.

The Managed Object pattern offers a number of benefits, but also some liabilities. On the positive side, the MIB and the Managed Object instances offer a single interface. This interface encapsulates the proprietary protocols of the Resource and simplifies the management of the system. A managing application is then able to communicate with each and every Resource in the MIB, regardless of its nature or particular characteristics. On the downside, the Managed Object provides only indirect access to the Resources, thus adding another level of indirection to the processing of management operations. This might compromise the system's performance, unless care is taken to keep the overhead to a minimum.

Note that the Managed Object must offer an interface that can gracefully support a large variety of proprietary interfaces. It must allow for the access to internal parameters and support the application of specific functions to the Resource. The Specification should express the details of the MIB in a manner as neutral and as precise as possible. In this way, the interoperability of the resource entities is ensured regardless of their underlying source or implementation.

### Remote Operation Pattern

In a distributed system such as a client/server system, the client of an operation is often removed from the location where that operation's implementation actually resides. The client must then access the implementation through a communication network. The problem here is to make a remote operation invocation appear exactly the same as a local operation invocation, both to the client and to the implementation. When attempting to solve this problem, the developer must deal with a number of conflicting forces. For one thing, making a call across the network is inherently more complex and less reliable than making a call simply across a processor's memory. Furthermore, both the client and server the application must be shielded from all network-specific details.

The proposed solution is to encapsulate all network interactions in stub objects, both on the Invoker's (client) and the Performer's (server) side. The Client Stub and the Server Stub communicate with each other using Connection and Message instances that are specific to the network under consideration. The Invoker and the Performer interact locally with their respective Stubs (see Figure 5). The result is a system where the invocation of a remote operation is decoupled from the network interactions needed to carry it out. This solution, referred to as Remote Operation pattern, can be seen as a refinement of the Proxy pattern [3]. It further decouples the network interactions from the Invoker and the Performer of the operation.

The Remote Operation pattern brings with it a number of benefits, but also some liabilities. On the positive side, the Invoker and the Performer are shielded from the network. Neither knows really if the invocation actually took place across the network or not. The Invoker doesn't even know where the operation is actually performed. To the Invoker, the whole business is no different from invoking a method on a local object, in this case the Client Stub.

On the negative side, a remote invocation takes longer than a local invocation. The time it takes to fulfill a remote invocation depends on the amount of overhead in the Messages and the bandwidth of the Connection. Furthermore, network errors may cause Messages to be lost. Some error correction mechanism is thus required to make the whole design more reliable. Also, it is hard to pass object references across process boundaries. The Client Stub and the Server Stub must handle pointers as well as object instances in a way that the semantics of the operation are preserved. All this extra activity adds to the time required to process a single remote operation.

## Discussion

Layla's pattern system illustrates the application and refinement of the pattern catalog presented in [3] for the NMI domain. Beyond this domain, the Layla patterns are relevant in the realm of distributed systems in general [2]. The three key patterns discussed here have been used in various contexts. The Manager-Agent pattern can be found in both CMIS and SNMP, and has been implemented in the OSIMIS [8] framework. The Managed Object pattern is used in CMIS as well, whereas in SNMP it occurs but in a diluted form. The Remote Operation pattern, finally, is being used in a number of interprocess protocols, such as CMIP and the protocol suite for X.400 email. Other implementations include Sun's and HP's remote procedure call libraries. Future OO frameworks in the NMI domain such as JMAPI [6] are likely to leverage these patterns, in case they implement CMIS or SNMP.

The current version of Layla is the result of several iterations, which seems to be quite typical for pattern-based framework development. Our pattern mining was definitely influenced by the NMI domain in that the numerous standards and tools of this domain inspired, and sometimes impeded, our work. And we can but confirm the lessons learned from framework development based on design patterns as reported elsewhere, for instance in [5, 11].

Layla is conceived as a pattern-based application framework. Its development demonstrates that pattern-based frameworks can be built for the demanding

NMI domain. Experimentation with Layla makes us believe that pattern benefits such as the flexibility, reusability, and documentation value of the framework and the resulting applications can indeed be reaped. Our experience suggests the pattern-based architecture of Layla makes NMI development considerably easier. We contend the pattern system upon which Layla is built will be helpful for other NMI framework builders and for NMI application developers alike. **C**

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